ROBUST SPR SYNTHESIS FOR LOW-ORDER POLYNOMIAL SEGMENTS AND INTERVAL POLYNOMIALS

Long Wang¹⁾ and Wensheng $Yu^{2)}$

 ¹⁾ Dept. of Mechanics and Engineering Science, Center for Systems and Control, Peking University, Beijing 100871, China, E-mail: longwang@mech.pku.edu.cn
 ²⁾ Lab of Engineering Science for Complex Systems, Institute of Automation, Chinese Academy of Sciences, Beijing 100080, China, E-mail: yws@compsys.ia.ac.cn

Abstract: We prove that, for low-order $(n \le 4)$ stable polynomial segments or interval polynomials, there always exists a fixed polynomial such that their ratio is SPR-invariant, thereby providing a rigorous proof of Anderson's claim on SPR synthesis for the fourth-order stable interval polynomials. Moreover, the relationship between SPR synthesis for low-order polynomial segments and SPR synthesis for low-order interval polynomials is also discussed.

Keywords: Synthesis Method, Strict Positive Realness(SPR), Constructive Design, Robust Stability, SPR-invariance, Transfer Functions, Polynomial Segments, Interval Polynomials, Polytopic Polynomials.

1. INTRODUCTION

The notion of strict positive realness (SPR) of transfer functions plays an important role in absolute stability theory, adaptive control and system identification[1-5]. In recent years, stimulated by the robustness analysis method[6-9], the study of robust strictly positive real transfer functions has received much attention, and great progress has been made[10-23]. However, most available results belong to the category of robust SPR analysis. Much work remains to be done in robust SPR synthesis.

Synthesis problems are mathematically more difficult than analysis problems. Usually, the synthesis problems require answering questions of existence and construction, whereas the analysis problems can be dealt with under the assumption of existence. Synthesis problems are of more practical significance from the engineering application viewpoint.

The basic statement of the robust strictly positive real synthesis problem is as follows: Given an *n*-th order robustly stable polynomial set F, does there exist, and how to construct a (fixed) polynomial b(s) such that, $\forall a(s) \in F$, a(s)/b(s) is strictly positive real? (If such a polynomial b(s) exists, then we say that F is synthesizable.)

When F is a low-order $(n \leq 3)$ interval polynomial set, the synthesis problem above has been considered by a number of authors and several important results^[13,14,16,17,19-21] have been presented. But when

F is a high-order $(n \ge 4)$ interval polynomial set, even in the case of n = 4, the synthesis problem above is still open^[16,17,19-21].

By the definition of SPR, it is easy to know that the Hurwitz stability of F is a necessary condition for the existence of polynomial b(s). In [13-15], it was proved that, if all polynomials in F have the same even (or odd) parts, such a polynomial b(s) always exists; In [13,14,16,19-21], it was proved that, if $n \leq 3$ and F is a stable interval polynomial set, such a polynomial b(s) always exists; Recent results in [18-20] show that, if n < 3 and F is the stable convex combination of two polynomials $a_1(s)$ and $a_2(s)$, such a polynomial b(s) always exists. Some sufficient condition for robust SPR synthesis are presented in [10,17,19-21], especially, the design method in [19,20] is numerically efficient for high-order polynomial segments and interval polynomials, and the derived conditions in [19,20] are necessary and sufficient for robust SPR synthesis of low-order $(n \leq 3)$ polynomial segments or interval polynomials.

It should be pointed out that. Anderson et al. [16] transformed the robust SPR synthesis problem for the fourth-order interval polynomial set into linear programming problem in 1990 (namely, equations (58)-(60) in [16]), and by using linear programming techniques, they concluded that such a linear programming problem always had a solution, thus, it was thought that the robust SPR synthesis problem for the fourth-order interval polynomial set had been solved. But in 1993, a synthesizable example in [17] showed that the corresponding linear programming problem had no solution. Hence, for the fourth-order interval polynomial set, on one hand, we could not prove theoretically the existence of robust SPR synthesis, on the other hand, we could not find a counterexample that is not synthesizable. Therefore, the robust SPR synthesis problem for interval polynomial set, even in the case of n = 4, is still an open problem^[16,17,13,14,19-21].

In this paper, we prove that, for low-order $(n \leq 4)$ stable polynomial segments or interval polynomials, there always exists a fixed polynomial such that their ratio is SPR-invariant, thereby providing a rigorous proof of Anderson's claim on SPR synthesis for the fourth-order sta-

ble interval polynomials. Moreover, the relationship between SPR synthesis for low-order polynomial segments and SPR synthesis for low-order interval polynomials is also discussed. Our proof is constructive, and is useful in solving the general SPR synthesis problem.

2. MAIN RESULTS

In this paper, P^n stands for the set of *n*-th order polynomials with real coefficients, R stands for the field of real numbers, $\partial(p)$ stands for the order of polynomial $p(\cdot)$, and $H^n \subset P^n$ stands for the set of *n*-th order Hurwitz stable polynomials.

In the sequel, $p(\cdot) \in P^m$, $q(\cdot) \in P^n$, f(s) = p(s)/q(s) is a rational function.

Definition 1^[10,16,17,23] A biproper rational function f(s) (i.e., $\partial(p) = \partial(q)$) is said to be strictly positive real(SPR), if

(i) f(s) is analytic in $\operatorname{Re}[s] \geq 0$, i.e., $q(\cdot) \in H^n$;

(ii) $\operatorname{Re}[f(j\omega)] > 0$, $\forall \omega \in \mathbb{R}$.

If f(s) = p(s)/q(s) is proper, it is easy to get the following property:

Lemma 1^[11] If f(s) = p(s)/q(s) is a proper rational function, $q(s) \in H^n$, and $\forall \omega \in R, \operatorname{Re}[f(j\omega)] > 0$, then $p(s) \in H^n \cup H^{n-1}$.

Denote $F = \{a_i(s) = s^n + \sum_{l=1}^n a_l^{(i)} s^{n-l}, i = 1, 2\}$ as the two endpoint polynomials of a stable polynomial segment \overline{F} (convex combination), it is easy to prove that:

Lemma 2^[16] $\forall a(s) \in \overline{F}, b(s)/a(s)$ is strictly positive real, if and only if, $b(s)/a_i(s), i = 1, 2$, are strictly positive real.

Consider an interval polynomials

$$K = \{a(s) = s^{n} + \sum_{i=1}^{n} a_{i}s^{n-i}, a_{i} \in [a_{i}^{-}, a_{i}^{+}], i = 1, 2, \cdots, n\}$$

Denote $F = \{a_i(s) = s^n + \sum_{l=1}^n a_l^{(i)} s^{n-l}, i = 1, 2, 3, 4\}$ as the four Kharitonov vertex polynomials of $K^{[6-9]}$.

Lemma 3^[6] K is robustly stable if and only if $a_i(s) \in H^n, i = 1, 2, 3, 4$.

The following result was proved by Dasgupta and Bhagwat^[10]:

Lemma 4^[10] $\forall a(s) \in K$, b(s)/a(s) is strictly positive real, if and only if, $b(s)/a_i(s)$, i = 1, 2, 3, 4, are strictly positive real.

First, for a low-order $(n \leq 3)$ stable convex combination of polynomials, by [18-20], we have

Theorem 1^[18] If $F = \{a_i(s) = s^n + \sum_{l=1}^n a_l^{(i)} s^{n-l}, i = 1, 2.\}$ is the set of the two endpoint polynomials of a low order $(n \leq 3)$ stable segment of polynomials (convex combination) \overline{F} , then there always exists a fixed polynomial b(s) such that $\forall a(s) \in \overline{F}, b(s)/a(s)$ is strictly positive real.

Furthermore, if F is the four Kharitonov vertex polynomials of a low-order $(n \leq 3)$ stable interval polynomial set, then we have

Theorem 2^[13,14,16,17,19-21] If $F = \{a_i(s) = s^n + \sum_{l=1}^n a_l^{(i)} s^{n-l}, i = 1, 2, 3, 4.\}$ is the set of the four Kharitonov vertex polynomials of a low order $(n \le 3)$ stable interval polynomial family K, then there always exists a fixed polynomial b(s) such that $\forall a(s) \in K, b(s)/a(s)$ is strictly positive real.

The following two theorems are the main results of this paper:

Theorem 3 If $F = \{a(s) = s^4 + a_1s^3 + a_2s^2 + a_3s + a_4, b(s) = s^4 + b_1s^3 + b_2s^2 + b_3s + b_4\}$ is the set of the two endpoint polynomials of a fourth order stable segment of polynomials (convex combination), then there always exists a fixed polynomial c(s) such that c(s)/a(s) and c(s)/b(s) are strictly positive real.

Consider the fourth-order interval polynomials

$$K = \left\{ \begin{array}{c} a(s) = s^4 + a_1 s^3 + a_2 s^2 + a_3 s + a_4, \\ a_i \in [a_i^-, a_i^+], i = 1, 2, 3, 4 \end{array} \right\}$$

Denote

$$\begin{aligned} a_1(s) &= s^4 + a_1^+ s^3 + a_2^+ s^2 + a_3^- s + a_4^- \\ a_2(s) &= s^4 + a_1^- s^3 + a_2^- s^2 + a_3^+ s + a_4^+ \\ a_3(s) &= s^4 + a_1^+ s^3 + a_2^- s^2 + a_3^- s + a_4^+ \\ a_4(s) &= s^4 + a_1^- s^3 + a_2^+ s^2 + a_3^+ s + a_4^- \end{aligned}$$

as the four Kharitonov vertex polynomials of $K^{[6-9]}$.

Theorem 4 If $F = \{a_i(s), i = 1, 2, 3, 4\}$ is the set of the four Kharitonov vertex polynomials of a fourth order stable interval polynomial family, then there always exists a fixed polynomial b(s) such that $\forall a(s) \in F$, b(s)/a(s) is strictly positive real.

Note that in Theorem 3, c(s)/a(s) and c(s)/b(s) being strictly positive real implies $\forall \lambda \in [0,1]$, $\frac{c(s)}{\lambda a(s)+(1-\lambda)b(s)}$ being strictly positive real (by Lemma 2); similarly, in Theorem 4, $\forall a(s) \in F$, b(s)/a(s) being strictly positive real implies $\forall a(s) \in K$, b(s)/a(s) being strictly positive real (by Lemma 4).

3. PROOFS OF MAIN RESULTS

In order to prove the main results above, we must first establish some lemmas.

Lemma 5 Suppose $a(s) = s^4 + a_1s^3 + a_2s^2 + a_3s + a_4 \in H^4$, then the following quadratic curve is an ellipse in the first quadrant of the x-y plane:

$$(a_2^2 - 4a_4)x^2 + 2(2a_3 - a_1a_2)xy + a_1^2y^2$$
$$-2(a_2a_3 - 2a_1a_4)x - 2a_1a_3y + a_3^2 = 0$$

and this ellipse is tangent with y axis at $(0, \frac{a_3}{a_1})$, tangent with the lines $x = a_1$ and $a_3y - a_4x = 0$ at $(a_1, a_2 - \frac{a_3}{a_1})$

and
$$(\frac{a_3}{a_2a_3-a_1a_4}, \frac{a_3a_4}{a_2a_3-a_1a_4})$$
, respectively.

Proof: Since a(s) is Hurwitz stable, Lemma 5 can be verified by a direct calculation.

Let $a(s) = s^4 + a_1s^3 + a_2s^2 + a_3s + a_4 \in H^4$, for notational simplicity, denote

apparently, Ω^a is a bounded convex set in the x-y plane.

Lemma 6 Suppose $a(s) = s^4 + a_1s^3 + a_2s^2 + a_3s + a_4 \in H^4$ and $(x, y) \in \Omega^a$, let $c(s) := s^3 + xs^2 + ys + \varepsilon$, where ε is positive and sufficiently small, then $\forall \omega \in R$, $\operatorname{Re}[\frac{c(j\omega)}{a(j\omega)}] > 0$.

Proof: Suppose $(x, y) \in \Omega^a$, let $c(s) := s^3 + xs^2 + ys + \varepsilon$, where $\varepsilon > 0$ and is sufficiently small.

 $\forall \omega \in R$, consider

$$\begin{array}{ll} \operatorname{Re}[\frac{c(j\omega)}{a(j\omega)}] &= \frac{1}{\mid a(j\omega)\mid^2}[(a_1-x)\omega^6\\ &+(a_2x-a_1y-a_3)\omega^4+(a_3y-a_4x)\omega^2\\ &+\varepsilon(\omega^4-a_2\omega^2+a_4)] \end{array}$$

In order to prove that $\forall \omega \in R$, $\operatorname{Re}[\frac{c(j\omega)}{a(j\omega)}] > 0$, let $t = \omega^2$, we only need to prove that, for any sufficiently small $\varepsilon > 0$,

$$\begin{array}{rl} f(t):=&t[(a_1-x)t^2+(a_2x-a_1y-a_3)t+(a_3y-a_4x)]\\ &+\varepsilon(t^2-a_2t+a_4)>0,\forall t\in[0,+\infty). \end{array}$$

Since $(x, y) \in \Omega^a$, by definition of Ω^a and Lemma 5, (x, y) satisfies $a_1 - x > 0$, $a_3y - a_4x > 0$, and

$$[a_2x - a_1y - a_3]^2 - 4(a_1 - x)(a_3y - a_4x) < 0$$

or

$$a_1 - x \ge 0, a_2x - a_1y - a_3 \ge 0, a_3y - a_4x > 0$$

Thus, $\forall t \in [0, +\infty)$

$$(a_1 - x)t^2 + (a_2x - a_1y - a_3)t + (a_3y - a_4x) > 0.$$

On the other hand, we have f(0) > 0, and for any $\varepsilon > 0$, if t is a sufficiently large or sufficiently small positive number, we have f(t) > 0. Namely, there exist $0 < t_1 < t_2$ such that, for all $\varepsilon > 0, t \in [0, t_1] \cup [t_2, +\infty), f(t) > 0$.

Denote

$$M = \inf_{t \in [t_1, t_2]} t[(a_1 - x)t^2 + (a_2x - a_1y - a_3)t + (a_3y - a_4x)],$$

$$N = \sup_{t \in [t_1, t_2]} t[t^2 - a_2t + a_3]$$

$$w = \sup_{t \in [t_1, t_2]} |t - u_2t + u_4|$$

then M > 0 and N > 0. Choosing $0 < \varepsilon < \frac{M}{N}$, by a direct calculation, we have

$$f(t) = t[(a_1 - x)t^2 + (a_2x - a_1y - a_3)t + (a_3y - a_4x)] + \varepsilon(t^2 - a_2t + a_4) > 0, \forall t \in [0, +\infty).$$

Namely

$$orall \omega \in R, \operatorname{Re}[rac{b(j\omega)}{a(j\omega)}] > 0$$

This completes the proof.

Lemma 7 Suppose $a(s) = s^4 + a_1s^3 + a_2s^2 + a_3s + a_4 \in H^4, b(s) = s^4 + b_1s^3 + b_2s^2 + b_3s + a_4 \in H^4$, if $\lambda b(s) + (1 - \lambda)a(s) \in H^4, \lambda \in [0, 1]$, then $\Omega_e^a \cap \Omega_e^b \neq \phi$.

Proof: If $\forall \lambda \in [0, 1]$, $\lambda b(s) + (1 - \lambda)a(s) \in H^4$, by Lemma 5, for any $\lambda \in [0, 1]$,

$$\begin{array}{rcl} \Omega_{e^{\lambda}}^{a_{\lambda}} := & \{(x,y) | (a_{\lambda 2}^2 - 4a_{\lambda 4})x^2 + 2(2a_{\lambda 3} - a_{\lambda 1}a_{\lambda 2})xy \\ & + a_{\lambda 1}^2 y^2 - 2(a_{\lambda 2}a_{\lambda 3} - 2a_{\lambda 1}a_{\lambda 4})x \\ & - 2a_{\lambda 1}a_{\lambda 3}y + a_{\lambda 3}^2 < 0 \} \end{array}$$

is also an elliptic region in the first quadrant of the x-y plane, where $a_{\lambda i} := a_i + \lambda (b_i - a_i), i = 1, 2, 3, 4$. Apparently, when λ changes continuously from 0 to 1, $\Omega_e^{a_\lambda}$ will change continuously from Ω_e^a to Ω_e^b .

Now assume $\Omega_e^{a} \cap \Omega_e^{b} = \phi$, by Lemma 5 (without loss of generality, suppose $\frac{b_3}{b_1} > \frac{a_3}{a_1}$), $\exists v \in [\frac{a_3}{a_1}, \frac{b_3}{b_1}]$ and $u \neq 0$, such that the following line l

$$l: \quad \frac{x}{u} + \frac{y}{v} = 1 \tag{1}$$

is tangent with Ω_e^a and Ω_e^b simultaneously, and l separates Ω_e^a and Ω_e^b (i.e., Ω_e^a and Ω_e^b are on different sides of l).

Since l is tangent with Ω_e^a , consider

$$\begin{cases} \frac{x}{u} + \frac{y}{v} = 1\\ (a_2^2 - 4a_4)x^2 + 2(2a_3 - a_1a_2)xy + a_1^2y^2\\ -2(a_2a_3 - 2a_1a_4)x - 2a_1a_3y + a_3^2 = 0 \end{cases}$$
(2)

since a(s) is Hurwitz stable and $u \neq 0$, by a direct calculation, we know that the necessary and sufficient condition for *l* being tangent with Ω_e^a is

$$uv^2 - a_1v^2 - a_2uv + a_3v + a_4u = 0$$
 (3)

Since l is tangent with Ω_e^b , for the same reason, we have

$$uv^2 - b_1v^2 - b_2uv + b_3v + b_4u = 0 \tag{4}$$

From (3) and (4), we obviously have $\forall \lambda \in [0, 1]$,

$$uv^2 - a_{\lambda 1}v^2 - a_{\lambda 2}uv + a_{\lambda 3}v + a_{\lambda 4}u = 0 \qquad (5)$$

(5) shows that l is also tangent with $\Omega_e^{\alpha_{\lambda}}(\forall \lambda \in [0, 1])$, but l separates Ω_e^{α} and Ω_e^{β} , and when λ changes continuously from 0 to 1, $\Omega_e^{\alpha_{\lambda}}$ will change continuously from Ω_e^{α} to Ω_e^{β} , which is obviously impossible. This completes the proof.

Lemma 8 If $F = \{a_i(s), i = 1, 2, 3, 4.\}$ is the set of the four Kharitonov vertex polynomials of a fourth order stable interval polynomial family, then $\Omega^{a_2} \subset \Omega^{a_4}$ and $\Omega^{a_3} \subset \Omega^{a_1}$.

Proof: By the definition of the notation Ω^a , it is easy to see that

Obviously, we have $\Omega^{a_2} \subset \Omega^{a_4}$ and $\Omega^{a_3} \subset \Omega^{a_1}$. This completes the proof.

Lemma 9 If $F = \{a_i(s), i = 1, 2, 3, 4\}$ is the set of the four Kharitonov vertex polynomials of a fourth order stable interval polynomial family, then $\bigcap_{i=1}^{4} \Omega^{a_i} \neq \phi$.

Lemma 9 plays an important role in proving Anderson's claim on robust SPR synthesis for the fourth-order stable interval polynomials. For a complete understanding of it, we give three different proofs in the sequel.

Proof 1: By Lemma 8, we only need to prove that $\Omega^{a_2} \cap \Omega^{a_3} \neq \phi$. By Lemma 7, we know that $\Omega^{a_2} \cap \Omega^{a_3}_e \neq \phi$, but $\Omega^{a_2} = \Omega^{a_2}_e \cup \Omega^{a_2}_t$ and $\Omega^{a_3} = \Omega^{a_3}_e \cup \Omega^{a_3}_t$, thus $\Omega^{a_2} \cap \Omega^{a_3} \neq \phi$. This completes the proof.

Proof 2: Since F is the set of the four Kharitonov vertex polynomials of a fourth order stable interval polynomial family, by Lemma 5, in the x-y plane, $\Omega_e^{a_2}$ and $\Omega_e^{a_4}$ are both tangent with x = 0 at $(0, \frac{a_3}{a_1^-})$ (denote this tangent point as A_{24}); $\Omega_e^{a_1}$ and $\Omega_e^{a_3}$ are both tangent with x = 0 at $(0, \frac{a_3}{a_1^+})$ (denote this tangent point as A_{13}). Denote the tangent point of $\Omega_e^{a_2}$ ($\Omega_e^{a_4}$) and $x = a_1^-$ as $A_2(a_1^-, a_2^- - \frac{a_3^+}{a_1^-})$ ($A_4(a_1^-, a_2^+ - \frac{a_3^+}{a_1^-})$); and denote the tangent point of $\Omega_e^{a_1}$ ($\Omega_e^{a_3}$) and $x = a_1^+$ as $A_1(a_1^+, a_2^+ - \frac{a_3^-}{a_1^+})$ ($A_3(a_1^+, a_2^- - \frac{a_3^-}{a_1^+})$). Furthermore, denote the intersection points of $x = a_1^-$ and the straight line $a_3^+ y - a_4^+ x =$ $0, a_3^+ y - a_4^- x = 0$ as $B_2(a_1^-, \frac{a_1^- a_4^+}{a_3^+}), B_4(a_1^-, \frac{a_1^- a_4^-}{a_3^+})$, respectively; and denote the intersection points of $x = a_1^+$ and the straight lines $a_3^- y - a_4^- x = 0, a_3^- y - a_4^+ x = 0$ as $B_1(a_1^+, \frac{a_1^+ a_4^-}{a_3^-}), B_3(a_1^+, \frac{a_1^+ a_4^+}{a_3^-})$, respectively.

In what follows, (A, B) stands for the set of points in the line segment connecting the point A and the point B, not including the endpoints A and B, [A, B)stands for the set of points in the line segment connecting the point A and the point B, including the endpoint A, but not B, (A, B] stands for the set of points in the line segment connecting the point A and the point B, including the endpoint B, but not A. Then it is easy to see that $[A_2, B_2) \subset \Omega^{a_2}, [A_2, B_2) \subset$ $[A_4, B_4) \subset \Omega^{a_4}, [A_3, B_3) \subset \Omega^{a_3}, [A_3, B_3) \subset [A_1, B_1) \subset$ Ω^{a_1} , and $(A_{24}, A_2] \subset \Omega^{a_2}, (A_{24}, A_2] \subset \Omega^{a_4}, (A_{13}, A_3] \subset \Omega^{a_5}, (A_{13}, A_3] \subset \Omega^{a_1}$.

Denote
$$A_3^{\star}$$
 as $(a_1^-, (\frac{a_2^-}{a_1^+} - 2\frac{a_3^-}{a_1^{+2}})a_1^- + \frac{a_3^-}{a_1^+})$, then $A_3^{\star} \in (A_{13}, A_3]$.

If $\frac{a_3^+}{a_1^-} = \frac{a_3^-}{a_1^+}$, i.e., $a_1^- = a_1^+$ and $a_3^- = a_3^+$. Then, take $\delta > 0, \delta$ sufficiently small, by Lemma 5, it is easy to verify that $(\delta, \frac{a_3^+}{a_1^-}) \in \cap_{i=1}^4 \Omega_e^{a_i}$, thus $\cap_{i=1}^4 \Omega^{a_i} \neq \phi$.

Now, suppose $\frac{a_3^+}{a_1^-} > \frac{a_3^-}{a_1^+}$ and $a_2^- - \frac{a_3^+}{a_1^-} \ge (\frac{a_2^-}{a_1^+} - 2\frac{a_3^-}{a_1^{+2}})a_1^- + \frac{a_3^-}{a_1^+}$

It is easy to verify that

$$(\frac{a_{2}^{-}}{a_{1}^{+}} - 2\frac{a_{3}^{-}}{a_{1}^{+2}})a_{1}^{-} + \frac{a_{3}^{-}}{a_{1}^{+}} > \frac{a_{1}^{-}a_{4}^{+}}{a_{3}^{+}}$$

Thus, we have $A_3^{\star} \in [A_2, B_2)$. Hence $A_3^{\star} \in [A_2, B_2) \cap (A_{13}, A_3]$. Therefore $A_3^{\star} \in \bigcap_{i=1}^4 \Omega^{a_i}$. Thus $\bigcap_{i=1}^4 \Omega^{a_i} \neq \phi$. Finally, with $\frac{a_3^+}{a_1^-} > \frac{a_3^-}{a_1^+}$, if

$$a_{2}^{-} - \frac{a_{3}^{+}}{a_{1}^{-}} < (\frac{a_{2}^{-}}{a_{1}^{+}} - 2\frac{a_{3}^{-}}{a_{1}^{+}})a_{1}^{-} + \frac{a_{3}^{-}}{a_{1}^{+}}$$

then it is easy to see that $(A_{13}, A_3] \cap (A_{24}, A_2] \neq \phi$ and $(A_{13}, A_3] \cap (A_{24}, A_2] \subset \bigcap_{i=1}^4 \Omega^{a_i}$. Thus, we also have $\bigcap_{i=1}^4 \Omega^{a_i} \neq \phi$. This completes the proof.

Proof 3:
$$A_{13}(0, \frac{a_3^-}{a_1^+}), A_{24}(0, \frac{a_3^+}{a_1^-}), B_2(a_1^-, \frac{a_1^-a_4^+}{a_3^+})$$
 and

 $B_3(a_1^+, \frac{a_1^+a_4^+}{a_3^-})$ are defined identically as in the Proof 2 above. (A, B) stands for the set of points in the line segment connecting the point A and the point B, but not including the endpoints A and B.

If $\frac{a_3^+}{a_1^-} = \frac{a_3^-}{a_1^+}$, i.e., $a_1^- = a_1^+$ and $a_3^- = a_3^+$. Then, take $\delta > 0, \delta$ sufficiently small, by Lemma 5, it is easy to verify that $(\delta, \frac{a_3^+}{a_1^-}) \in \cap_{i=1}^4 \Omega_e^{a_i}$, thus $\cap_{i=1}^4 \Omega^{a_i} \neq \phi$.

Now, suppose $\frac{a_3^+}{a_1^-} > \frac{a_3^-}{a_1^+}$, then it is easy to see that $(A_{13}, B_3) \cap (A_{24}, B_2) \neq \phi$ and $(A_{13}, B_3) \cap (A_{24}, B_2) \subset \cap_{i=1}^4 \Omega^{a_i}$. Thus, we also have $\cap_{i=1}^4 \Omega^{a_i} \neq \phi$. This completes the proof.

Lemma 10 Suppose $a(s) = s^4 + a_1s^3 + a_2s^2 + a_3s + a_4 \in$ $H^4, b(s) = s^3 + xs^2 + ys + z$, and $\forall \omega \in R$, $\operatorname{Re}[\frac{b(j\omega)}{a(j\omega)}] > 0$, take

$$\widetilde{b}(s) := b(s) + r \cdot c(s), \quad r > 0, r \text{ sufficiently small}$$

where c(s) is a fixed fourth-order monic polynomial. Then $\frac{\widetilde{b}(s)}{a(s)}$ is strictly positive real.

Proof: Obviously, $\partial(a) = \partial(\widetilde{b})$, namely, $\widetilde{b}(s)$ and a(s) have the same order. Since $a(s) \in H^4$, there exists $\omega_1 > 0$ such that, for all $|\omega| \ge \omega_1$, $\operatorname{Re}[\frac{\widetilde{b}(j\omega)}{a(j\omega)}] > 0$. Denote

$$M_1 = \inf_{|\omega| \le \omega_1} \operatorname{Re}[rac{b(j\omega)}{a(j\omega)}] \qquad N_1 = \sup_{|\omega| \le \omega_1} |\operatorname{Re}[rac{c(j\omega)}{a(j\omega)}]|$$

Then $M_1 > 0$ and $N_1 > 0$. Choosing $0 < r < \frac{M_1}{N_1}$, it can be directly verified that

$$orall \omega \in R, {
m Re}[rac{\widetilde{b}(j\omega)}{a(j\omega)}] > 0$$

This completes the proof.

Now Theorem 3 is proved by simply combining Lemmas 5-7 and Lemma 10. Theorem 4 is proved by simply combining Lemmas 5-6 and Lemmas 9-10.

4. DISCUSSIONS AND EXAMPLES

The following three examples correspond to different cases in the proof of our main results.

Example 1 Suppose $a_1(s) = s^4 + 89s^3 + 56s^2 + 88s + 1$, $a_2(s) = s^4 + 11s^3 + 56s^2 + 88s + 50$, $a_3(s) = s^4 + 89s^3 + 56s^2 + 88s + 50$, $a_4(s) = s^4 + 11s^3 + 56s^2 + 88s + 1$ are the four Kharitonov vertex polynomials of a fourth-order interval polynomial set K, it is easy to check using Kharitonov's Theorem that K is robustly stable. By our method as in the constructive proof of Theorem 4, it is easy to get $(11, 7.6657) \in \bigcap_{i=1}^4 \Omega^{a_i}$. Thus, choose $b(s) = s^3 + 11s^2 + 7.76657s + \varepsilon$, where ε is a sufficiently small positive number (ε is determined by Lemma 6, in this example, $0 < \varepsilon \leq 3$), take $\varepsilon = 2$, by Lemma 6, $\forall \omega \in R$, $\operatorname{Re}[\frac{b(j\omega)}{a_i(j\omega)}] > 0$, i = 1, 2, 3, 4. Finally, let $\tilde{b}(s) := b(s) + r \cdot s^4$, where r > 0, r.

sufficiently small (r is determined by Lemma 10, in this example, $0 < r \le 0.5$), it is easy to check that $\frac{\widetilde{b}(s)}{a_i(s)}$, i =

1, 2, 3, 4, are strictly positive real (note that b(s) and $\tilde{b}(s)$ are not unique).

In this example, if we take $F = \{a_1(s) = s^4 + 11s^3 + 56s^2 + 88s + 50, a_2(s) = s^4 + 89s^3 + 56s^2 + 88s + 50\}$, then it is exactly the counter-example provided in [17]. It can be checked that F does not satisfy the sufficient conditions in [10,16,17], but we can use the methods in [13-15,19,20] to do SPR synthesis. When F is enlarged to the interval polynomial set K in this example, the synthesis methods in [13-15] fail too, but we can still use the methods in [19,20] to do synthesis. It is quite straightforward to do synthesis using the method in this paper.

Example 2 Suppose $a_1(s) = s^4 + 5s^3 + 6s^2 + 4s + 0.5, a_2(s) = s^4 + 2s^3 + 6s^2 + 6s + 1, a_3(s) = s^4 + 5s^3 + 6s^2 + 4s + 1, a_4(s) = s^4 + 2s^3 + 6s^2 + 6s + 0.5$ are the four Kharitonov vertex polynomials of a fourth-order interval polynomial set K, it is easy to check using Kharitonov's Theorem that K is robustly stable. By our method as in the constructive proof of Theorem 4, it is easy to get $(2, 2.56) \in \cap_{i=1}^4 \Omega^{a_i}$. Thus, choose $b(s) = s^3 + 2s^2 + 2.56s + \varepsilon$, where ε is a sufficiently small positive number (in this example, $0 < \varepsilon \le 1$), take $\varepsilon = 0.5$, by Lemma 6, $\forall \omega \in R$, $\operatorname{Re}[\frac{b(j\omega)}{a_i(j\omega)}] > 0, i = 1, 2, 3, 4$. Finally, let $\widetilde{b}(s) := b(s) + r \cdot s^4$, where r > 0, r sufficiently small (in this example,

 $0 < r \le 0.5$), it is easy to check that $\frac{\widetilde{b}(s)}{a_i(s)}, i = 1, 2, 3, 4$, are strictly positive real.

Example 3 Suppose $a_1(s) = s^4 + 2.5s^3 + 6s^2 + 4s + 0.5, a_2(s) = s^4 + 2s^3 + 5s^2 + 6s + 5, a_3(s) = s^4 + 2.5s^3 + 5s^2 + 4s + 5, a_4(s) = s^4 + 2s^3 + 6s^2 + 6s + 0.5$ are the four Kharitonov vertex polynomials of a fourth-order interval polynomial set K, it is easy to check using Kharitonov's Theorem that K is robustly stable. By our method as in the constructive proof of Theorem 4, it is easy to get $(1.1475, 2.4262) \in \bigcap_{i=1}^4 \Omega^{a_i}$. Thus, choose $b(s) = s^3 + 1.1475s^2 + 2.4262s + \epsilon$, where ϵ is a sufficiently small positive number (in this example, $0 < \epsilon \le 1$), take $\epsilon = 0.5$, by Lemma 6, $\forall \omega \in R$, $\operatorname{Re}[\frac{b(j\omega)}{a_i(j\omega)}] > 0, i =$

1,2,3,4. Finally, let $\tilde{b}(s) := b(s) + r \cdot s^4$, where r > 0, r sufficiently small (in this example, $0 < r \le 0.2$), it is easy $\tilde{b}(s)$

to check that $\frac{\widetilde{b}(s)}{a_i(s)}$, i = 1, 2, 3, 4, are strictly positive real.

Remark 1 From the proofs of Theorem 3 and Theorem 4, we can see that, this paper not only proves the existence, but also provides a design procedure.

Remark 2 Lemma 10 actually holds for arbitrary n-th order polynomials^[19,20].

Remark 3 The constructive synthesis method is also insightful and helpful in solving the general robust SPR synthesis problem. In fact, we have recently succeeded in proving the existence on robust SPR synthesis for fifth-order stable convex combinations using a similar method^[19]. The SPR synthesis for higher-order systems is currently under investigation.

Remark 4 Robust stability of a polynomial segment can be checked by many efficient methods, e.g., eigenvalue method, root locus method, value set method, etc.^[8,9]. Robust stability of K in Theorem 4 can be ascertained by checking only two Kharitonov vertex polynomials^[24].

Remark 5 From the proofs of Lemma 9, we can establish the relationship between SPR synthesis for the fourth-order polynomial segments and SPR synthesis for the fourth-order interval polynomials. In fact, it is easy to see that Theorem 3 implies Theorem 4. Similarly, Theorem 1 implies Theorem 2. However, similar results may not be true for higher-order $(n \ge 5)$ systems. This subject is currently under investigation.

Remark 6 Our results can easily be generalized to discrete-time case.

Finally, it should also be pointed out that, for the vertex set $F = \{a_i(s) = s^n + \sum_{l=1}^n a_l^{(i)} s^{n-l}, i = 1, 2, \cdots, m.\}$ of a general polytopic polynomial family \overline{F} , even if \overline{F} is robustly stable, it is still possible that there does not exist a polynomial $c(s) \in H^{n-1}$, such that, $\forall \omega \in R$, $\operatorname{Re}[\frac{c(j\omega)}{a(j\omega)}] >$

0, for all $a(s) \in \overline{F}$.

To see this, let us look at an example of a third order triangle polynomial family.

Example 4 Let $F = \{a_1(s) = s^3 + 2.6s^2 + 37s + 64, a_2(s) = s^3 + 17s^2 + 83s + 978, a_3(s) = s^3 + 15s^2 + 64s^3 + 15s^2 + 15s^$

28s + 415}. It is easy to verify that $a_i(s), i = 1, 2, 3$, are Hurwitz stable. Moreover, all edges of \overline{F} , i.e., $\lambda a_i(s) + (1 - \lambda)a_j(s), \lambda \in [0, 1], i, j = 1, 2, 3$, are also Hurwitz stable. Therefore, by Edge Theorem^[6-9], \overline{F} is robustly stable. On the other hand, by a direct computation, we can easily see that there does not exist a polynomial $c(s) \in H^2$, such

that $\forall \omega \in R$, $\operatorname{Re}[\frac{c(j\omega)}{a_i(j\omega)}] > 0, i = 1, 2, 3$. Note that, in this example, though there does not exist

Note that, in this example, though there does not exist a polynomial $c(s) \in H^2$ such that $\forall \omega \in R$, $\operatorname{Re}[\frac{c(j\omega)}{a_i(j\omega)}] > 0, i = 1, 2, 3$. But if we take $\tilde{c}(s) = s^3 + 6s^2 + 73s + 68$, it is easy to check $\frac{\tilde{c}(s)}{a_i(s)}, i = 1, 2, 3$, are strictly positive real. This shows some intrinsic differences between the SPR synthesis of interval polynomial families and the SPR synthesis of polytopic polynomial families. This problem deserves further investigation.

5. CONCLUSIONS

We have proved that, for low-order $(n \leq 4)$ stable polynomial segments or interval polynomials, there always exists a fixed polynomial such that their ratio is SPR-invariant, thereby providing a rigorous proof of Anderson's claim on SPR synthesis for the fourth-order stable interval polynomials. Moreover, the relationship between SPR synthesis for low-order polynomials has also been discussed.

Acknowledgments

This work was supported by the National Key Project of China, the National Natural Science Foundation of China (69925307), Natural Science Foundation of Chinese Academy of Sciences and National Lab of Intelligent Control and Systems of Tsinghua University.

References

- Kalman, R. E. (1963), Lyapunov functions for the problem of Lur'e in automatic control. Proc. Nat. Acad. Sci.(USA), 49: 201-205.
- [2] Popov, V. M. (1973), Hyperstability of Automatic Control Systems. New York: Springer-Verlag.
- [3] Desoer, C. A. and Vidyasagar, M. (1975), Feedback Systems: Input-Output Properties. San Diego: Academic Press.
- [4] Anderson, B. D. O. and Moore, J. B. (1970), Linear Optimal Control. New York: Prentice Hall.
- [5] Landau, Y. D. (1979), Adaptive Control: The Model Reference Approach. New York: Marcel Dekker.
- [6] Kharitonov, V. L. (1978), Asymptotic stability of an equilibrium position of a family of systems of linear differential equations. Differentsial'nye Uravneniya, 14: 2086-2088.
- [7] Bartlett, A. C., Hollot, C. V., and Huang, L. (1988), Root locations for an entire polytope of polynomial: it suffices to check the edges. *Math. Contr. Signals Syst.*, 1: 61-71.

- [8] Bhattacharyya, S. P., Chapellat H., and Keel, L. H. (1995), Robust Control - The Parametric Approach. New York: Prentice Hall.
- [9] Barmish, B. R. (1994), New Tools for Robustness of Linear Systems. New York: MacMillan Publishing Company.
- [10] Dasgupta, S. and Bhagwat, A. S. (1987), Conditions for designing strictly positive real transfer functions for adaptive output error identification. *IEEE Trans. Circuits Syst.*, CAS-34: 731-737.
- [11] Chapellat, H., Dahleh, M., and Bhattacharyya, S. P. (1991), On robust nonlinear stability of interval control systems. *IEEE Trans. Automat. Contr.*, AC-36: 59-69.
- [12] Wang, L. and Huang, L. (1991), Finite Verification of Strict Positive Realness of Interval Rational Functions. *Chinese Science Bulletin*, 36: 262-264.
- [13] Hollot, C. V., Huang, L., and Xu, Z. L. (1989), Designing strictly positive real transfer function families: A necessary and sufficient condition for low degree and structured families. Proc. of Mathematical Theory of Networks and Systems (eds. Kaashoek, M. A., Van Schuppen, J. H., Ran, A. C. M.). Boston, Basel, Berlin: Birkh äuser, 215-227.
- [14] Huang, L, Hollot, C. V., and Xu, Z. L. (1990), Robust analysis of strictly positive real function set. Preprints of The Second Japan-China Joint Symposium on Systems Control Theory and its Applications. 210-220.
- [15] Patel, V. V. and Datta, K. B. (1997), Classification of units in H_{∞} and an alternative proof of Kharitonov's theorem. *IEEE Trans. Circuits Syst.; Part I,* CAS-44: 454-458.
- [16] Anderson, B. D. O., Dasgupta, S., Khargonekar, P., et al. (1990), Robust strict positive realness: characterization and construction. *IEEE Trans. Circuits Syst.*, CAS-37: 869-876.
- [17] Betser, A. and Zeheb, E. (1993), Design of robust strictly positive real transfer functions. *IEEE Trans. Circuits* Syst.; Part I, CAS-40: 573-580.
- [18] Yu, W. S. and Huang, L. (1999), A necessary and sufficient conditions on robust SPR stabilization for low degree systems. *Chinese Science Bulletin*, 44: 517-520.
- [19] Wang, L. and Yu, W. S. (2000), Complete characterization of strictly positive real regions and robust strictly positive real synthesis method. *Science in China*, E-43: 97-112.
- [20] Wang, L. and Yu, W. S. (1999), A new approach to robust synthesis of strictly positive real transfer functions. Stability and Control: Theory and Applications, 2: 13-24.
- [21] Marquez, H. J. and Agathoklis, P. (1998), On the existence of robust strictly positive real rational functions. *IEEE Trans. Circuits Syst.; Part I, CAS-45:* 962-967.
- [22] Yu, W. S., Wang, L., and Tan, M. (1999), Complete characterization of strictly positive real regions in coefficient space. Proceedings of the IEEE Hong Kong Symposium on Robotics and Control, Hong Kong, 259-264.
- [23] Yu, W. S. and Wang, L. (1999), Some remarks on the definitions of strict positive realness. Proc. of Chinese Conference on Decision and Control, Northeast University Press, Shenyang, 135-139.
- [24] Anderson, B. D. O., Jury, E., and Mansour, M. (1987), On robust Hurwitz polynomials. *IEEE Trans. on Automatic Control*, AC-32: 909-913.
- [25] Wang, L. and Yu, W. S. (2000), On robust stability of polynomials and robust strict positive realness of transfer functions, *IEEE Trans. Circuits Syst.*; Part I, (to appear).
- [26] Yu, W. S. and Wang, L., Robust SPR Synthesis for Fifth-Order Convex Combinations, Technical Report, Chinese Academy of Sciences, 2000.